QIP Course 9: Quantum Factorization Algorithm (Part 2)

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Inverse QFT

Answers to the previous exercises will be given on the blackboard.

Let $|0\rangle$, ..., $|N-1\rangle$ be an orthonormal basis of an N-dimensional space. The QFT transforms

$$|j\rangle \mapsto \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} \exp(2\pi i j k/N) |k\rangle.$$

The inverse of QFT (IQFT) is given by

$$|k\rangle \mapsto \frac{1}{\sqrt{N}} \sum_{\ell=0}^{N-1} \exp(-2\pi i k\ell/N) |\ell\rangle.$$
 (1)

IQFT can be realized by applying R_k^{-1} and H^{-1} in the reverse order. \Rightarrow IQFT can also be realized with the same efficiency (n(n+1)/2) operations of R_k^{-1} and H^{-1}) as QFT.

Phase estimation 1

Suppose that we have a unitary matrix U and its eigenvector vector $|u\rangle$. Let $\exp(2\pi i\theta)$ be the eigenvalue to which $|u\rangle$ belongs to. We shall show how we can compute θ .

Assumption: We are able to do the controlled- U^{2^j} operation for any $j \ge 0$. Suppose that we apply the controlled- U^{2^j} to $(|0\rangle + |1\rangle)|u\rangle$, with $|u\rangle$ being the target (we omit the normalizing factor $1/\sqrt{2}$). Then the result is

$$|0\rangle|u\rangle + |1\rangle \otimes U^{2^{j}}|u\rangle$$

$$= |0\rangle|u\rangle + |1\rangle \otimes \exp(2\pi i 2^{j}\theta)|u\rangle$$

$$= (|0\rangle + \exp(2\pi i 2^{j}\theta)|1\rangle) \otimes |u\rangle$$

Assume we have t qubits that are initialized to $(|0\rangle + |1\rangle)/\sqrt{2}$, and apply the controlled- U^{2^j} to the j-th qubit (the rightmost is the zero-th). The result is

$$\frac{1}{2^{t/2}}(|0\rangle + \exp(2\pi i 2^{t-1}\theta)|1\rangle) \otimes \cdots \otimes (|0\rangle + \exp(2\pi i 2^{0}\theta)|1\rangle)$$

$$= \frac{1}{2^{t/2}} \sum_{k=0}^{2^{t-1}} \exp(2\pi i k\theta)|k\rangle. \tag{2}$$

Applying the IQFT (1) to (2) yields

$$\frac{1}{2^t} \sum_{\ell=0}^{2^t-1} \sum_{k=0}^{2^t-1} \exp\left(\frac{-2\pi i k \ell}{2^t}\right) \exp(2\pi i k \theta) |\ell\rangle.$$

Probability distribution of the measurement outcomes 1

$$\frac{1}{2^t} \sum_{\ell=0}^{2^t-1} \sum_{k=0}^{2^t-1} \exp\left(\frac{-2\pi i k \ell}{2^t}\right) \exp(2\pi i k \theta) |\ell\rangle.$$

We shall compute the probability distribution of the mesurement in the $\{|0\rangle, |1\rangle, |2\rangle, ..., |2^t - 1\rangle\}$ basis. (The observable is $\sum_{j=0}^{2^t-1} j|j\rangle\langle j|$.) Recall that $0 \le \theta < 1$, and we can write

$$\theta = 0.b_1b_2\cdots b_tb_{t+1}\cdots.$$

Let $b = b_1 b_2 \cdots b_t$. We have $0 \le b \le 2^t - 1$. b is the nearest t-bit integer $\le 2^t \theta$. When m is the measurement outcome, we regard $m/2^t$ as our estimate of θ . I will explain that $m \simeq 2^t \theta \simeq 2^t b$ with large probability.

$$\frac{1}{2^t} \sum_{\ell=0}^{2^t-1} \sum_{k=0}^{2^t-1} \exp\left(\frac{-2\pi i k \ell}{2^t}\right) \exp(2\pi i k \theta) |\ell\rangle.$$

Let α_c be the coefficient of $|(b+c) \mod 2^t\rangle$ in the result of the IQFT (the above). We shall show that if c is large then $|\alpha_c|$ is small. Observe that the coefficient of $|\ell\rangle$ is

$$\frac{1}{2^t} \sum_{k=0}^{2^t - 1} \exp\left(\frac{-2\pi i k \ell}{2^t}\right) \exp(2\pi i k \theta) = \frac{1}{2^t} \sum_{k=0}^{2^t - 1} \left[\exp\left(2\pi i (\theta - \ell/2^t)\right)\right]^k$$

Substituting ℓ with b + c we have

$$\alpha_c = \frac{1}{2^t} \sum_{k=0}^{2^t - 1} [\exp(2\pi i(\theta - (b+c)/2^t))]^k$$

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is the sum of a geometric series, so it is equal to

$$\alpha_c = \frac{1}{2^t} \cdot \frac{1 - \exp(2\pi i (2^t \theta - (b+c)))}{1 - \exp(2\pi i (\theta - (b+c)/2^t))}.$$

Define $\delta = \theta - b/2^t$, then

$$\alpha_c = \frac{1}{2^t} \cdot \frac{1 - \exp(2\pi i (2^t \delta - c))}{1 - \exp(2\pi i (\delta - c/2^t))}$$

We shall upper bound the probability of getting a measurement outcome m such that |m - b| > e. Observe $\Pr[m = b + c] = |\alpha_c|^2$.

We shall upper bound the probability of getting a measurement outcome m such that |m-b| > e. We have

$$p(|m-b| > e) = \sum_{-2^{t-1} < c \le -e-1, e+1 \le c < 2^{t-1}} |\alpha_c|^2.$$

Since $|1 - \exp(ix)| \le 2$,

$$|\alpha_c| \le \frac{2}{2^t |1 - \exp(2\pi i (\delta - c/2^t))|}.$$

We have $|1 - \exp(ix)| \ge 2|x|/\pi$ for $-\pi \le x \le \pi$ and $-\pi \le 2\pi(\delta - c/2^t) \le \pi$, it follows

$$|\alpha_c| \le \frac{1}{2^{t+1}|\delta - c/2^t|}.$$

Therefore we have

$$\begin{split} 4p(|m-b| > e) & \leq & \sum_{-2^{t-1} < c \leq -e-1} \frac{1}{(2^t \delta - c)^2} + \sum_{e+1 \leq c < 2^{t-1}} \frac{1}{(2^t \delta - c)^2} \\ & \leq & \sum_{-2^{t-1} < c \leq -e-1} \frac{1}{c^2} + \sum_{e+1 \leq c < 2^{t-1}} \frac{1}{(c-1)^2} \\ & \leq & 2 \sum_{e \leq c < 2^{t-1}-1} \frac{1}{c^2} \\ & \leq & 2 \int_{e-1}^{2^{t-1}-1} \frac{dc}{c^2} \\ & \leq & 2 \int_{e-1}^{\infty} \frac{dc}{c^2} \\ & = & \frac{2}{(e-1)}. \end{split}$$

Sufficiently many qubits ensure the accuracy with high probability

Suppose that we want an accuracy of 2^{-n} , that is, $|\theta - m/2^t| < 2^{-n}$.

$$\begin{aligned} |\theta - m/2^t| &< 2^{-n} \\ \Leftrightarrow & |2^t \theta - m| &< 2^{t-n} \\ \Leftrightarrow & |b - m| &< 2^{t-n} - 1. \end{aligned}$$

We can see that $e = 2^{t-n} - 1$ ensures the desired accuracy. The probability of the accuracy below 2^{-n} is $1/2(2^{t-n} - 2)$. In order for $1/2(2^{t-n} - 2) < \epsilon$, we need $t \ge n + \log_2(2 + 1/2\epsilon)$.

Exercise

1. Let

$$U = \left(\begin{array}{cc} 1 & 0\\ 0 & \exp(2\pi i 5/16) \end{array}\right)$$

Find the all eigenvalues of *U*.

- 2. Let $|u\rangle$ be the eigenvector of U and assume $U|u\rangle \neq |u\rangle$. Assume that we do the phase estimation with t=3. Then there is eight possible measurement outcomes. Compute the probability distribution of outcomes **and their corresponding estimates of** θ **.** I recommend you to use Mathematica, Matlab, Maple, and so on.
- 3. By using $p(|m-b| > e) \le \frac{1}{2(e-1)}$ compute the lower bound on the probability of the event that the mesurement outcome of θ is within 3/8 from the true value $\theta = 5/16$. How much difference exists between the lower bound and the true probability?