Fundamentals of Mathematical and Computing Sciences: Applied Mathematical Sciences

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Outline of the Lecture

This course introduces several basic concepts of mathematical optimization, probability and statistics, and is intended to provide key knowledge necessary for advanced study in Mathematical and Computing Sciences.

Evaluation

Average of three reports.

Part I: Mathematical Optimization

References

[Ben-Tal-Nemirovski] A. Ben-Tal and A. Nemirovski, *Lectures on Modern Convex Optimization: Analysis, Algorithms, and Engineering Applications* (SIAM, Philadelphia, PA, 2001).

[Lasserre] J. B. Lasserre, *Moments, Positive Polynomial and their Applications* (Imperial College Press, London, 2010).

[Renegar] J. Renegar, *A Mathematical View of Interior-Point Methods in Convex Optimization* (SIAM, Philadelphia, PA, 2001).

[Schrijver] A. Schrijver, *Theory of Linear and Integer Programming* (John Wiley and Sons, Chichester, 1986).

1 Preliminaries

We assume that convexity and closeness (openness) of sets are familiar concepts for the readers.

Let $\langle\ ,\ \rangle$ be an arbitrary inner product on $\mathbb{R}^n$. Given a linear operator $\mathcal{A} : \mathbb{R}^n \longrightarrow \mathbb{R}^m$, its *adjoint* operator $\mathcal{A}^* : \mathbb{R}^m \longrightarrow \mathbb{R}^n$ is such that

$$\langle \mathcal{A}(\boldsymbol{x}), \boldsymbol{y} \rangle = \langle \boldsymbol{x}, \mathcal{A}^*(\boldsymbol{y}) \rangle, \quad \forall \boldsymbol{x} \in \mathbb{R}^n,\ \forall \boldsymbol{y} \in \mathbb{R}^m.$$

Definition 1.1 A set $\mathcal{K} \subseteq \mathbb{R}^n$ is called *cone* if for any positive scalar $\alpha > 0$ and an arbitrary element $\boldsymbol{x}$ of $\mathcal{K}$, $\alpha\boldsymbol{x} \in \mathcal{K}$.

Definition 1.2 A cone is said to be *pointed* if $\mathcal{K} \cap -\mathcal{K} = \{\mathbf{0}\}$.

Definition 1.3 Given a cone $\mathcal{K} \subseteq \mathbb{R}^n$, its *dual* cone is defined as $\mathcal{K}^* := \{\boldsymbol{x} \in \mathbb{R}^n \mid \langle \boldsymbol{x}, \boldsymbol{y} \rangle \geq 0,\ \forall \boldsymbol{y} \in \mathcal{K}\}$.

Definition 1.4 If a cone is such that $\mathcal{K}^* = \mathcal{K}$, it is called **self-dual**.

Theorem 1.5 (Separation theorem for convex sets [Ben-Tal-Nemirovski]) Let A, B nonempty non-intersecting convex subsets of $\mathbb{R}^n$. Then, $\exists \boldsymbol{s} \in \mathbb{R}^n$, $\boldsymbol{s} \neq 0$ such that

$$\sup_{\boldsymbol{a} \in A} \langle \boldsymbol{a}, \boldsymbol{s} \rangle \leq \inf_{\boldsymbol{b} \in B} \langle \boldsymbol{b}, \boldsymbol{s} \rangle.$$

2 Conic Linear Program

The Linear Program (LP) is the most basic mathematical optimization problem. We will start defining a generalization of the LP.

The *Conic Linear Program* (CLP) is defined as follows[1]:

$$\text{(CLP)} \left\{ \begin{array}{ll} \text{minimize} & \langle \boldsymbol{c}, \boldsymbol{x} \rangle \\ \text{subject to} & \mathcal{A}(\boldsymbol{x}) = \boldsymbol{b}, \\ & \boldsymbol{x} \in \mathcal{K}, \end{array} \right.$$

where $\boldsymbol{c} \in \mathbb{R}^n$, $\boldsymbol{b} \in \mathbb{R}^m$, $\mathcal{A}(\cdot)$ is an linear operator, and $\mathcal{K}$ is a <u>closed convex cone</u> in $\mathbb{R}^n$.

[1] strictly speaking, the term "minimize" should be replaced by "infimum"

The dual problem of (CLP) is defined as[2]:

$$\text{(DCLP)} \begin{cases} \text{maximize} & \langle \boldsymbol{b}, \boldsymbol{y} \rangle \\ \text{subject to} & \mathcal{A}^*(\boldsymbol{y}) + \boldsymbol{s} = \boldsymbol{c}, \\ & \boldsymbol{s} \in \mathcal{K}^*, \end{cases}$$

where the inner product is defined on $\mathbb{R}^m$ now. Notice that $\mathcal{K}^*$ is a closed convex cone, too.

Example 2.1 If we chose $\mathcal{K} = \mathbb{R}^n_+$, $\mathcal{A} := \boldsymbol{A} \in \mathbb{R}^{m \times n}$, and $\langle \boldsymbol{c}, \boldsymbol{x} \rangle = \boldsymbol{c}^T \boldsymbol{x}$, (CLP) becomes an LP. Likewise, taking $\mathcal{K} = \mathbb{S}^n_+$, the cone of positive semidefinite symmetric matrices, and the inner product which defines the Frobenius norm, we have a *Semidefinite Program* (SDP); $\mathcal{K} = \mathbb{Q}^n_+ := \{\boldsymbol{x} \in \mathbb{R}^n \mid x_1^2 \geq \sum_{i=2}^n x_i^2\}$, the second-order cone, we have a *Second-Order Cone Program* (SOCP).

The following result known as *weak duality* is a simple consequence of above facts.

Lemma 2.2 (Weak Duality) Let $\boldsymbol{x}$ be feasible for (CLP) and $(\boldsymbol{y}, \boldsymbol{s})$ feasible for (DCLP). Then $\langle \boldsymbol{b}, \boldsymbol{y} \rangle \leq \langle \boldsymbol{c}, \boldsymbol{x} \rangle$.

Proof:
$\langle \boldsymbol{c}, \boldsymbol{x} \rangle - \langle \boldsymbol{b}, \boldsymbol{y} \rangle = \langle \boldsymbol{x}, \boldsymbol{c} \rangle - \langle \mathcal{A}(\boldsymbol{x}), \boldsymbol{y} \rangle = \langle \boldsymbol{x}, \boldsymbol{c} - \mathcal{A}^*(\boldsymbol{y}) \rangle = \langle \boldsymbol{x}, \boldsymbol{s} \rangle \geq 0$ since $\boldsymbol{x} \in \mathcal{K}$ and $\boldsymbol{s} \in \mathcal{K}^*$. ∎

The following example shows that **strong duality** does not hold in general.

$$\begin{cases} \text{minimize} & \left\langle \begin{pmatrix} -1 \\ 0 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} \right\rangle \\ \text{subject to} & \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\ & \boldsymbol{x} \in \mathcal{K} = \{\boldsymbol{x} \in \mathbb{R}^4 \mid x_1^2 + x_2^2 \leq x_3^2,\ x_3, x_4 \geq 0\} \end{cases}$$

$$\begin{cases} \text{maximize} & \left\langle \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} \right\rangle \\ \text{subject to} & \begin{pmatrix} -1 \\ 0 \\ 0 \\ 0 \end{pmatrix} - \begin{pmatrix} y_1 \\ 0 \\ 0 \\ y_1 \end{pmatrix} - \begin{pmatrix} 0 \\ y_2 \\ y_2 \\ 0 \end{pmatrix} \in \mathcal{K}^* = \mathcal{K}. \end{cases}$$

Both problems are feasible, but the optimal value of the primal is 0 while for the dual is -1.

Theorem 2.3 (Strong Duality) If (CLP) is bounded from below and it is strictly feasible (*i.e*, $\exists \boldsymbol{x} \in \text{int}(\mathcal{K})$ and $\mathcal{A}(\boldsymbol{x}) = \boldsymbol{b}$), then (DCLP) is solvable and its optimal value coincides with the one of (CLP). The result is valid if the roles of (CLP) and (DCLP) are exchanged.

Proof:
Let c_{val} be the optimal value of (CLP) which exists by the assumption. We need to show that (DCLP) is solvable and have the same optimal value.

For $\boldsymbol{c} = \boldsymbol{0}$, $c_{\text{val}} = 0$ and the existence of the feasible solution $\boldsymbol{y} = \boldsymbol{0}$, $\boldsymbol{s} = \boldsymbol{0}$ for (DCLP) is evident.

Now let $\boldsymbol{c} \neq \boldsymbol{0}$. Consider the set

$$M = \{\boldsymbol{x} \in \mathbb{R}^n \mid \mathcal{A}(\boldsymbol{x}) = \boldsymbol{b},\ \langle \boldsymbol{c}, \boldsymbol{x} \rangle \leq c_{\text{val}}\}.$$

[2]strictly speaking, the term “maximize” should be replaced by “supremum”